ELFS, TREES AND QUANTUM WALKS



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ELFS, TREES AND QUANTUM WALKS

simple graph G = (V, E)

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unit electric flow f from source s to sink M



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effective resistance $R = \sum_{e} f_{e}^{2}$



1. sample edge *e* with probability $p_e = f_e^2/R$



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2. pick random endpoint *x* from *e*, set source s = x



process ends when step 2. picks sink vertex

• elfs

coupling with random walk

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coupling with random walk

• trees

electric hitting time is $O(\log n)$

• elfs

coupling with random walk

• trees

electric hitting time is $O(\log n)$

• quantum walks

quantum algorithm for elfs

elfs process describes Markov chain

$$\{Y_0 = s, Y_1, Y_2, \dots, Y_{\rho} \in M\}$$

with ρ the "electric hitting time"



random walk

$$\{X_0=s,X_1,X_2,\ldots,X_\tau\in M\}$$

with τ the random walk hitting time



lemma: exists stopping times $0 < \nu_1 < \cdots < \nu_{\rho} = \tau$ such that

$$\{Y_0 = X_0, Y_1 = X_{\nu_1}, \ldots, Y_{\rho} = X_{\nu_{\rho}} = X_{\tau}\}$$



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i.e., exists coupling between random walk and elfs process

lemma: exists stopping times $0 < \nu_1 < \cdots < \nu_{\rho} = \tau$ such that

$$\{Y_0 = X_0, Y_1 = X_{\nu_1}, \ldots, Y_{\rho} = X_{\nu_{\rho}} = X_{\tau}\}$$



corollary: random walk and elfs process same arrival distribution

$$\Pr(Y_{\rho} = x) = \Pr(X_{\tau} = x) = \sum_{y} f_{y,x}, \qquad x \in M$$

stopping rule ν from s: when visiting vertex x, stop with probability¹

$$p_{x} = \frac{\sum_{y} (v_{x} - v_{y})^{2}}{v_{x}d_{x} + \sum_{y} (v_{x} - v_{y})^{2}}$$

 v_z = voltage at *z* in unit electric *s*-*M* flow

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e.g.,
$$p_x = 1$$
 for $x \in M$

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expected stopping time $\mathbb{E}[\nu]$?

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lemma:
$$\mathbb{E}[\nu] = \frac{1}{R} \sum_{x} v_x^2 d_x =: \mathrm{ET}_s$$

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using
$$v_x \le v_s = R$$
: 1 $\le Rd_s \quad \left(= v_s d_s = 1/\Pr_s(\tau_M < \tau_s^+)\right)$
 $\le ET_s$
 $\le HT_s \quad \left(=\sum_x v_x d_x\right)$

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lemma': $\operatorname{ET}_s = \text{"escape time" from } s \text{ to } M$ = $\mathbb{E}\{1 + \max\{t \mid X_t = s\}\}$

 v_z = voltage at z in unit electric s-M flow

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 $\operatorname{EHT}_{s} \sim \min\{|M|, n/|M|\}$

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conquer via bottoms-up:
steps below vertex = f(# steps below children)

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quantum walks

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=

algorithmic tool in quantum computing (element distinctness, search, Monte Carlo)

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original motivation for elfs: quantum walks can create quantum flow state

$$\left|f\right\rangle = \frac{1}{\sqrt{R}}\sum_{e}f_{e}\left|e\right\rangle$$

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$$\left|f
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angle$$

measuring $|f\rangle$ yields electric flow sample

let
$$\operatorname{Ref}_{+} = \operatorname{reflection} \operatorname{around} \operatorname{span}_{x,y}\{|x, y\rangle + |y, x\rangle\}$$

 $\operatorname{Ref}_{*} = \operatorname{reflection} \operatorname{around} \operatorname{span}_{x}\left\{\sum_{y \sim x} |x, y\rangle\right\}$



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$$|x, y\rangle + |y, x\rangle$$
}
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quantum walk operator

 $W = \operatorname{Ref}_{+}\operatorname{Ref}_{*}$

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invariant subspaces:

$$\Pi_+ \cap \Pi_* = \sum_{x,y} |x,y\rangle$$

 $\Pi_{-} \cap (\Pi_{*})^{\perp} = \{ |h\rangle \mid h \text{ closed flow} \}$

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! electric flow state $|f\rangle$ orthogonal to all closed flows

quantum walk algorithm (for *s*-*t* flow) **1.** start from $|s, t\rangle$

2. apply "quantum phase estimation" to project orthogonal to invariant subspace

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(ignoring errors in remainder)

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returns element from sink *M* according to random walk arrival distribution

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many applications (to explore) (sampling RSTs, semi-supervised learning, linear system solving)

general framing: (quantum) algorithms for Lx=b

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quantum/probabilistic solution

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quantum/probabilistic solution

 $\begin{array}{c} \operatorname{quantum} \operatorname{in} \\ \operatorname{poly}(\kappa) \end{array}$ via HHL / quantum walks

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explicit solution

quantum in $\widetilde{O}(\sqrt{mn}/\varepsilon)$ via Grover [Apers-de Wolf, FOCS'20]

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 \rightarrow polynomial speedups

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no in general (BQP-complete) yes for low rank (Tang '19) explicit solution

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 \rightarrow exponential speedups?

explicit solution

quantum?

classical in n^{ω} via matrix inversion

 \rightarrow no speedups?

Elfs, trees and quantum walks Summary



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• electric flow sampling with random walks?

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 - explore quantum applications

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Thank you!